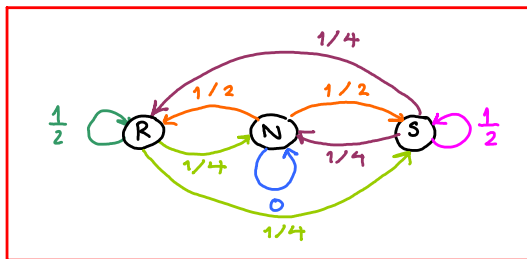


(a)

As hinted, we draw three states.

Next, follow the description sentence by sentence to get the transition probabilities.

- ① Never have two nice days in a row
- ② If have a nice day, just as likely to have snow as rain the next day
- ③ If have snow or rain, they have an even chance of having the same the next day.
- ④ If there is change from snow or rain, only half of the time is this a change to a nice day.

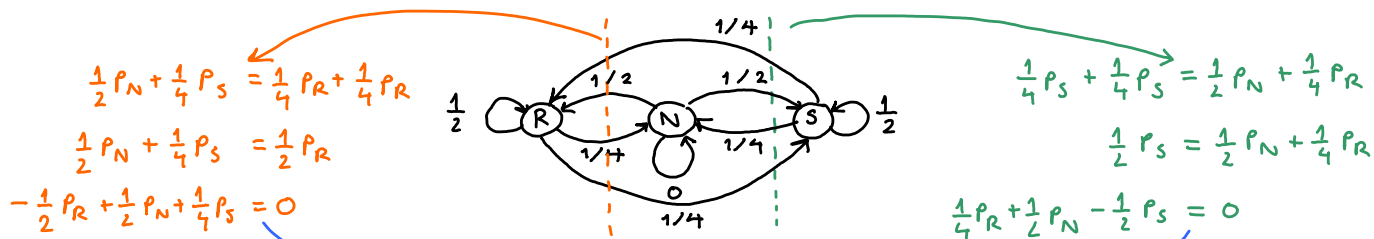


Note that we set the color of each arrow and its label to match with the corresponding description above.

(b) The corresponding transition probability matrix is

$$P = \begin{matrix} & \begin{matrix} R & N & S \end{matrix} \\ \begin{matrix} R \\ N \\ S \end{matrix} & \begin{bmatrix} 1/2 & 1/4 & 1/4 \\ 1/2 & 0 & 1/2 \\ 1/4 & 1/4 & 1/2 \end{bmatrix} \end{matrix}$$

(c) (i) We first write two balance equations from the two boundaries shown below:



One more equation: $p_R + p_N + p_S = 1$

Solve 3 eqns, 3 unknowns.

$$p_R = 0.4, p_N = 0.2, \text{ and } p_S = 0.4$$

(ii) The steady-state probabilities can also be found in MATLAB via eigen-decomposition.

First, we need to put the matrix P into MATLAB:

$$P = \text{sym}([1/2 \ 1/4 \ 1/4; \ 1/2 \ 0 \ 1/2; \ 1/4 \ 1/4 \ 1/2]);$$

Method 1:

In class, we have seen how to find the steady-state probabilities via finding the limiting distribution of

$$P^{(n)} = P^{(0)} P^n \quad (\text{Alternatively, we may write } P^{(n)} = P^{(1)} P^{(n-1)})$$

To find the limit of P^n , we use eigen-decomposition:

```
% Method 1
[V,D] = eig(P)
Dinf = (abs(D)==1)
Pinf = V*Dinf*(V^(-1))
p = Pinf(1,:)
eval(p)
```

```
V =
[ 1, -1, 1]
[ 1, 0, -4]
[ 1, 1, 1]
D =
[ 1, 0, 0]
[ 0, 1/4, 0]
[ 0, 0, -1/4]
Dinf =
1 0 0
0 0 0
0 0 0
Pinf =
[ 2/5, 1/5, 2/5]
[ 2/5, 1/5, 2/5]
[ 2/5, 1/5, 2/5]
p =
[ 2/5, 1/5, 2/5]
ans =
0.4000 0.2000 0.4000
```

$$P = V D V^{-1}$$

$$P^n = V D^n V^{-1}$$

$$\lim_{n \rightarrow \infty} P^n = V \left(\lim_{n \rightarrow \infty} D^n \right) V^{-1} \equiv P^\infty$$

$$D^\infty \equiv \lim_{n \rightarrow \infty} D^n$$

so, we have

$$\lim_{n \rightarrow \infty} P^{(n)} = P^{(0)} \begin{bmatrix} 2/5 & 1/5 & 2/5 \\ 2/5 & 1/5 & 2/5 \\ 2/5 & 1/5 & 2/5 \end{bmatrix}$$

$$= \begin{bmatrix} P_R^{(0)} & P_N^{(0)} & P_S^{(0)} \end{bmatrix} \begin{bmatrix} 2/5 & 1/5 & 2/5 \\ 2/5 & 1/5 & 2/5 \\ 2/5 & 1/5 & 2/5 \end{bmatrix} = \begin{bmatrix} 2/5 & 1/5 & 2/5 \end{bmatrix}$$

Method 2:

When the system has already reached its limiting distribution p , we expect that, in the next slot/step, the system will still have the same distribution. Now, going from one slot/step to the next is simply a multiplication by the transition probability matrix P . Therefore, we must have

$$p = p P.$$

Equivalently, $p^T = P^T p^T.$

Let's compare this with $\lambda \vec{v} = A \vec{v}$. We see that p^T is simply the eigen vector of $A = P^T$ whose corresponding eigen value is $\lambda = 1$:

Note the use of P^T

```
% Method 2
[V D] = eig(P');
k = find(diag(D)==1);
v = V(:,k);
p = (v./sum(v))';
eval(p)
```

find $\lambda = 1$
 find the corresponding eigen-vector
 Normalize the eigen vector to have $\sum p_i = 1$

```
p =
[ 2/5, 1/5, 2/5]
ans =
0.4000 0.2000 0.4000
```

(d)

```

n = 1e4;          % The number of slots to be considered
S = [1,2,3];     % Two possible states
P = [1/2 1/4 1/4; 1/2 0 1/2; 1/4 1/4 1/2]; % Transition probability matrix
X1 = 2;          % Initial state

```

```

X = MarkovChainGS(n,S,P,X1);

```

```

% Approximate the transition probabilities from the simulation

```

```

P_sim = []; x = X(1:(n-1)); y = X(2:n);

```

```

for k = 1:length(S)

```

```

    l = find(x==S(k)); LI = length(l);

```

```

    yc = y(l); cond_rel_freq = hist(yc,S)/LI;

```

```

    P_sim = [P_sim; cond_rel_freq];

```

```

end

```

```

P_sim

```

```

% Approximate the proportions of time that the states occur

```

```

p_sim = hist(X,S)./n

```

```

>> MarkovChain_2014_HW5_Q1_d

```

```

P_sim =

```

```

    0.4995    0.2553    0.2452

```

```

    0.5177         0    0.4823

```

```

    0.2583    0.2468    0.4949

```

```

p_sim =

```

```

    0.4090    0.2008    0.3902

```

Simulation results are close enough to the theoretical results obtained in part (c).

(e) Let today be slot $\neq 0$. Then $P^{(0)} = \begin{bmatrix} R & N & S \\ 0 & 0 & 1 \end{bmatrix}$.

$$(i) P^{(1)} = P^{(0)} P = \begin{bmatrix} 0 & 0 & 1 \\ 1/2 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/2 \end{bmatrix} = \begin{bmatrix} R & N & S \\ 1/4 & 1/4 & 1/2 \end{bmatrix}$$

$\hookrightarrow p_N^{(1)} = 1/4.$

$$(ii) P^{(2)} = P^{(1)} P = \begin{bmatrix} 1/4 & 1/4 & 1/2 \\ 1/2 & 1/4 & 1/4 \\ 1/4 & 1/4 & 1/2 \end{bmatrix} = \begin{bmatrix} 3/8 & 3/16 & 7/16 \end{bmatrix}$$

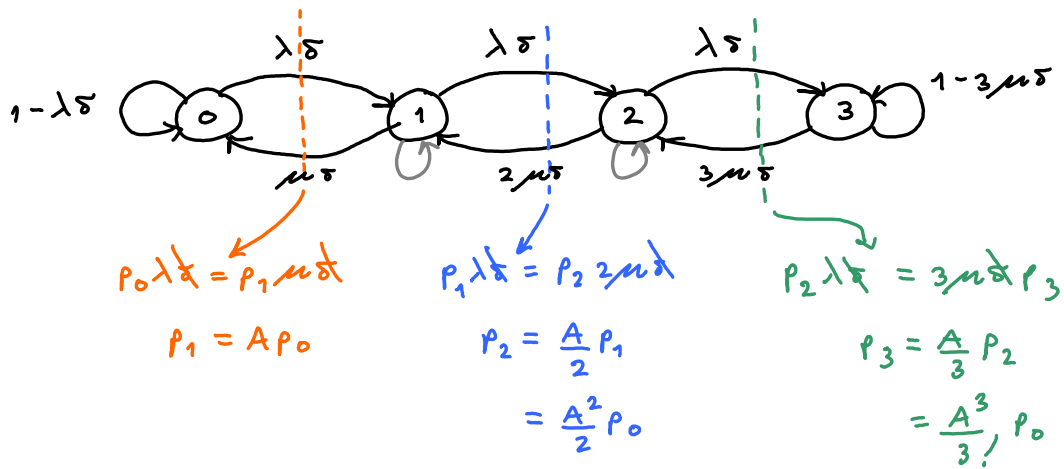
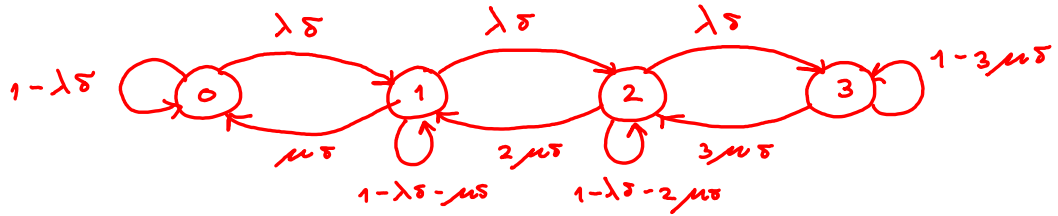
$\hookrightarrow p_N^{(2)} = 3/16.$

(iii) "365 days" is a long time. $P^{(365)} \approx P = \begin{bmatrix} 0.2 & 0.2 & 0.7 \end{bmatrix}$
 The probability of being a nice day $\approx p_N = 0.2.$

$m = 3$

(a) Erlang B model

Markov chain:



$$p_0 + p_1 + p_2 + p_3 = 1 \Rightarrow p_0 = \left(1 + A + \frac{A^2}{2} + \frac{A^3}{3!} \right)^{-1}$$

$$A = \frac{\lambda}{\mu} = \lambda \times \frac{1}{\mu} = \left(10 \frac{\text{calls}}{\text{hour}} \times \frac{1 \text{ hour}}{60 \text{ mins}} \right) \times (12 \text{ mins}) = 2 \text{ Erlangs.}$$

$$\Rightarrow p_0 = \frac{3}{19}, p_1 = \frac{6}{19}, p_2 = \frac{6}{19}, p_3 = \frac{4}{19}$$

↓

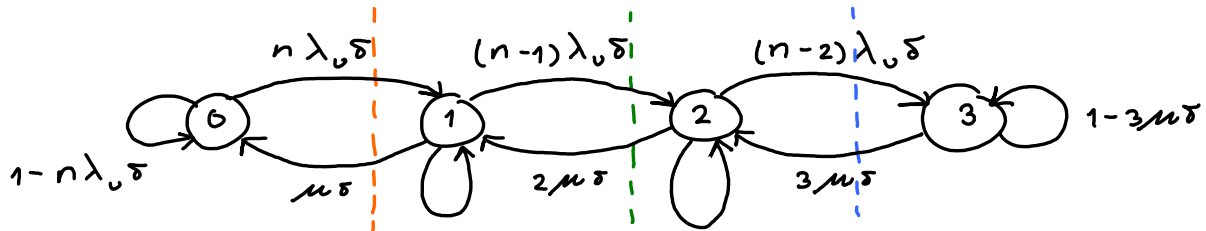
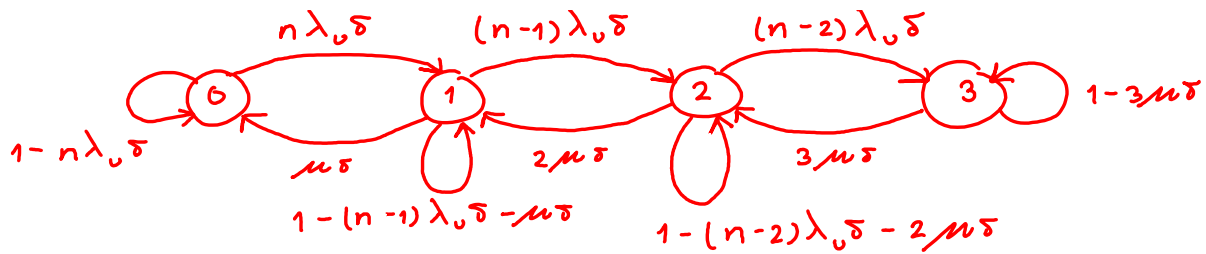
$$\text{call blocking probability} = \frac{4}{19} \approx 0.211$$

(b) and (c)

Observe that $\lambda_u \times n = \lambda$ in part (a).

so, $\lambda_u = \frac{\lambda}{n}$ and $A_u = \frac{A}{n} = \frac{2}{n}$ Erlangs.

Markov chain:



$$p_0 n \lambda_0 \delta = p_1 \mu \delta$$

$$p_1 = n A_0$$

$$p_1 (n-1) \lambda_0 \delta = 2 \mu \delta p_2$$

$$p_2 = \frac{(n-1)}{2} A_0 p_1$$

$$= \frac{n(n-1)}{2} A_0^2 p_0$$

$$p_2 (n-2) \lambda_0 \delta = p_3 3 \mu \delta$$

$$p_3 = \frac{(n-2)}{3} A_0 p_2$$

$$= \frac{n(n-1)(n-2)}{3!} A_0^3 p_0$$

$$p_0 + p_1 + p_2 + p_3 = 1 \Rightarrow \begin{cases} p_0 = \frac{25}{131}, p_1 = \frac{50}{131}, p_2 = \frac{40}{131}, p_3 = \frac{16}{131} & \text{when } n=5, \\ p_0 = 0.159, p_1 = 0.379, p_2 = 0.316, p_3 = 0.206 & \text{when } n=100. \end{cases}$$

As discussed in class, the call blocking probability is given by

$$\frac{(n-m) p_m}{\sum_{k=0}^m (n-k) p_k} = \frac{(n-3) p_3}{n p_0 + (n-1) p_1 + (n-2) p_2 + (n-3) p_3}$$

$$= \begin{cases} \frac{32}{477} \approx 0.067 & \text{when } n=5 \\ 0.203 & \text{when } n=100 \end{cases}$$

↑
close to the answer from Erlang B.

Remark:

For those who are interested in why the Engset model converges to the Erlang B model when $n \rightarrow \infty$, read on.

Note that

$$p_k = \binom{n}{k} A^k p_0 = \binom{n}{k} \frac{A^k}{n^k} p_0 = \frac{n!}{(n-k)! k! n^k} A^k p_0.$$

For fixed k ,

$$\frac{n!}{(n-k)! n^k} = \frac{n \times (n-1) \times \dots \times (n-(k-1))}{n^k} = \frac{n}{n} \times \frac{n-1}{n} \times \dots \times \frac{n-(k-1)}{n}$$

$$\rightarrow 1 \text{ as } n \rightarrow \infty.$$

Hence,

$$p_k = \frac{\binom{n}{k} \frac{A^k}{n^k}}{\sum_{i=0}^m \binom{n}{i} \frac{A^i}{n^i}} \rightarrow \frac{\frac{1}{k!} A^k}{\sum_{i=0}^m \frac{1}{i!} A^i} \text{ as } n \rightarrow \infty$$

↑
same as the steady-state probabilities in Erlang B model.

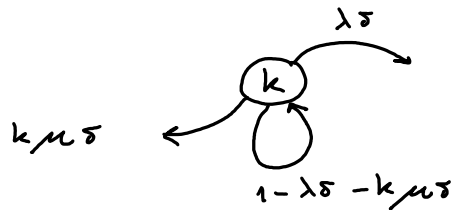
Similarly, for the call blocking probability,

$$p_{CB} = \frac{(n-m) p_m}{\sum_{k=0}^m (n-k) p_k} = \frac{(n-m) \binom{n}{m} \frac{A^m}{n^m} p_0}{\sum_{k=0}^m (n-k) \binom{n}{k} \frac{A^k}{n^k} p_0} = \frac{\binom{n}{m} \frac{A^m}{n^m}}{\sum_{k=0}^m \frac{n-k}{n-m} \binom{n}{k} \frac{A^k}{n^k}}$$

$$\rightarrow \frac{\frac{A^m}{m!}}{\sum_{k=0}^m \frac{A^k}{k!}} \text{ as } n \rightarrow \infty$$

↪ same as the call blocking probability in Erlang B model.

(a) Nothing changes from the $M/M/m/m$ model when $k < m$.
 We have

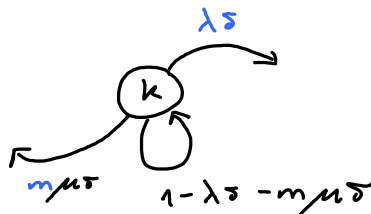


When $k \geq m$, the call request rate is still λ . The difference is that now we have a queue for the new requests to wait. (In $M/M/m/m$, these requests are discarded and the calls are blocked.)

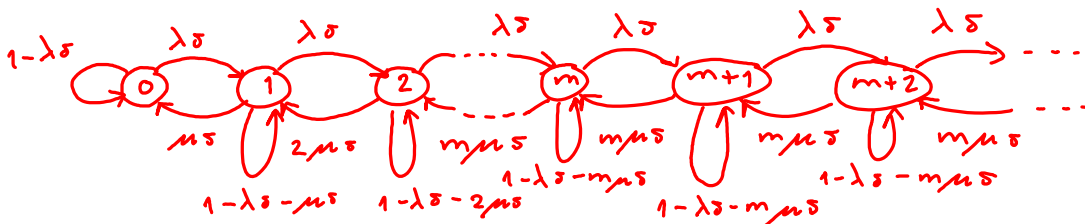
When $k \geq m$, all m channels are being used. There are $k - m$ requests waiting in the queue. When there is one new call request, it will be added to the queue and hence the system move from state k to $k + 1$. Again, this new call request occurs with probability $\lambda \delta$ (approximately).

When $k \geq m$, all m channels are being used. There are m customers talking on the phone. So, the probability of one call ends is (approximately) $m \mu \delta$.

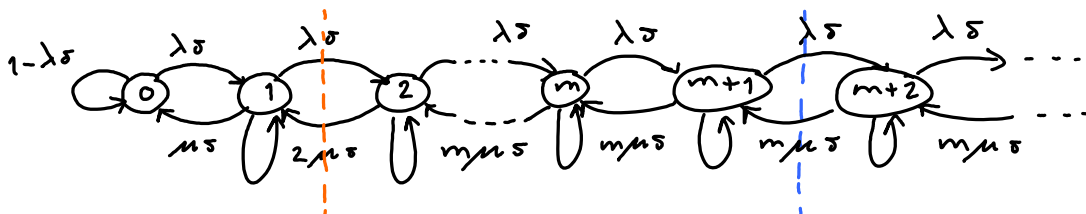
Therefore, when $k \geq m$, we have

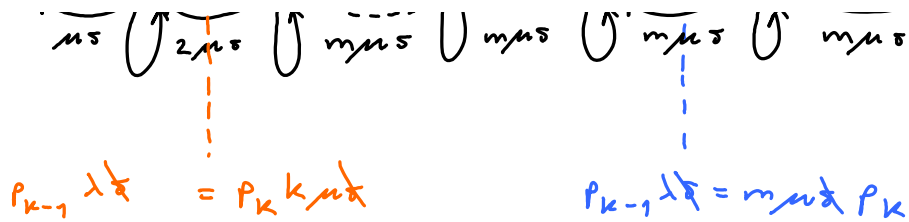


Markov chain :



(b)





$$p_k = \frac{A}{k} p_{k-1}$$

$$= \frac{A^k}{k!} p_0$$

for $0 < k < m$

$$p_k = \frac{A}{m} p_{k-1}$$

for $k \geq m$

$$\Rightarrow p_m = \frac{A}{m} p_{m-1}$$

$$= \frac{A}{m} \frac{A^{m-1}}{(m-1)!} p_0$$

$$= \frac{A^m}{m!} p_0$$

$$p_k = \left(\frac{A}{m}\right)^{k-m} \frac{A^m}{m!} p_0 = \frac{A^k}{m!(m^{k-m})} p_0$$

$$= \frac{m^m}{m!} \left(\frac{A}{m}\right)^k p_0 \quad \text{for } k \geq m.$$

$$\sum_{k=0}^m p_k = 1 \Rightarrow 1 = \sum_{k=0}^{m-1} \frac{A^k}{k!} p_0 + \sum_{k=m}^{\infty} \frac{m^m}{m!} \left(\frac{A}{m}\right)^k p_0 = \sum_{k=0}^{m-1} \frac{A^k}{k!} p_0 + \frac{A^m}{m!} \left(1 - \frac{A}{m}\right)^{-1} p_0$$

geometric series $\Rightarrow \begin{cases} \frac{m^m}{m!} \frac{\left(\frac{A}{m}\right)^m}{1 - \frac{A}{m}} p_0 & \text{if } A < m \\ \infty & \text{if } A \geq m \end{cases}$

$$\Rightarrow p_0 = \left(\left(\sum_{k=0}^{m-1} \frac{A^k}{k!} \right) + \frac{A^m}{m! \left(1 - \frac{A}{m}\right)} \right)^{-1}$$

Therefore,

$$p_k = \begin{cases} \frac{A^k}{k!} p_0, & k < m \\ \frac{A^k}{m!(m^{k-m})} p_0, & k \geq m \end{cases}$$

(c) Delayed call probability

$$= \sum_{k=m}^{\infty} p_k = \frac{A^m}{m!} \frac{1}{1 - A/m} p_0 = \frac{A^m}{m! (1 - A/m)}$$

$$\begin{aligned}
 &= \sum_{k=m}^{\infty} p_k = \frac{A^m}{m! \left(1 - \frac{A}{m}\right)} p_0 = \frac{m! \left(1 - \frac{A}{m}\right)}{\frac{A^m}{m! \left(1 - \frac{A}{m}\right)} + \sum_{k=0}^{m-1} \frac{A^k}{k!}} \\
 &= \frac{A^m}{A^m + m! \left(1 - \frac{A}{m}\right)} \sum_{k=0}^{m-1} \frac{A^k}{k!}
 \end{aligned}$$

Remark: This formula is call the "Erlang C formula".

(a)

$$P_m = \frac{\binom{n}{m} A_v^m}{\sum_{k=0}^m \binom{n}{k} A_v^k} = \frac{\sum_{k=0}^m \binom{n}{k} A_v^k - \sum_{k=0}^{m-1} \binom{n}{k} A_v^k}{\sum_{k=0}^m \binom{n}{k} A_v^k}$$

$$= \frac{z(m, n) - z(m-1, n)}{z(m, n)} = 1 - \frac{z(m-1, n)}{z(m, n)}$$

Hence, $c = 1$

(b)

First, note that

$$(n-k) \times \binom{n}{k} = (n-k) \times \frac{n!}{k!(n-k)!} = \frac{n!}{k!(n-k-1)!}$$

$$= \frac{n \times (n-1)!}{k!(n-1-k)!} = n \binom{n-1}{k}$$

Therefore,

$$P_b = \frac{(n-m) \binom{n}{m} A_v^m}{\sum_{k=0}^m (n-k) \binom{n}{k} A_v^k} = \frac{\cancel{n} \binom{n-1}{m} A_v^m}{\sum_{k=0}^m \cancel{n} \binom{n-1}{k} A_v^k}$$

$$= \frac{\sum_{k=0}^m \binom{n-1}{k} A_v^k - \sum_{k=0}^{m-1} \binom{n-1}{k} A_v^k}{\sum_{k=0}^m \binom{n-1}{k} A_v^k}$$

$$= \frac{z(m, n-1) - z(m-1, n-1)}{z(m, n-1)} = 1 - \frac{z(m-1, n-1)}{z(m, n-1)}$$

Hence, $c_1 = c_2 = c_4 = 1$ and $c_3 = 0$.

...

(c)

$$P_b = \frac{\binom{n-1}{m} A_u^m}{\sum_{k=0}^m \binom{n-1}{k} A_u^k} \stackrel{m=n-1}{=} \frac{\binom{m}{m} A_u^m}{\sum_{k=0}^m \binom{m}{k} A_u^k} = \frac{A_u^m}{(1+A_u)^m} = \left(\frac{A_u}{1+A_u} \right)^m$$